

A proof of Example 4.6 (p114) that more closely resembles the pumping lemma

Claim $L = \{a^k b^k : k \geq 0\}$ is not regular.

Proof: Assume L is regular and argue for a contradiction.

As L is regular, there is a dfa $M = \{Q, \{a, b\}, S, q_0, F\}$ accepting L . Set $m = |Q| + 1$, so m is one more than the number of states in M . Consider the sequence of m states defined by $\delta^*(q_0, a^i)$, $i = 1, 2, \dots, m$.

There are more elements in this sequence than states in the dfa, so at least one state must be repeated.

[This is called the pigeonhole principle].

Pick one of the repeated states and call it q .

Thus, for two distinct values r and s in $\{1, 2, \dots, m\}$

we have $\delta^*(q_0, a^r) = \delta^*(q_0, a^s) = q$.

We may assume $r < s$ (otherwise swap them).

So let $l = s-r$. l is the length of a cycle of 'a' links starting and ending at q . Formally,

$$\delta^*(q, a^{r+nl}) = q \text{ for any } n=0,1,2,\dots$$

We already know that M accepts $a^m b^m$, so

$$\delta^*(q_0, a^m b^m) = \hat{q} \text{ for some } \hat{q} \text{ in } F.$$

But we can insert a cycle of l a's after the first r a's without affecting the result!

Thus

$$\delta^*(q_0, a^{m+nl} b^m) = \hat{q} \text{ for } n=0,1,2,\dots$$

But $\hat{q} \in F$. So in other words, M accepts $a^{m+nl} b^m$

for $n=0,1,2,\dots$, contradicting the fact that M accepts only elements of L . □

Briefer version you can use in UT Quiz

Assume L is regular, and a dfa $M = \{Q, \Sigma, \delta, q_0, F\}$ accepts it. Let $m = |Q| + 1$ and consider $\delta^*(q_0, a^i)$, $i=1, \dots, m$.

By the pigeonhole principle, there exists a repeated state q in this sequence; i.e. we have r, s with $1 \leq r < s \leq m$ such that $\delta^*(q_0, a^r) = \delta^*(q_0, a^s) = q$.

Thus, a cycle of a -links of length $s-r$ starts at q ,

$$\text{so } \delta^*(q_0, a^m b^m) = \delta^*(q_0, a^{m+n(s-r)} b^m)$$

for $n = 0, 1, 2, \dots$

This contradicts the fact that M accepts only strings of the form $a^k b^k$.

□