

Strategies for basing the CS theory
course on non-decision problems...

... and using real computer programs.

i.e. a practical approach to teaching the
CS theory course

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What is the CS “theory course”?

Usually covers some or all of:

- automata theory
 - finite automata, context free languages, Turing machines
- computability theory
 - undecidability of halting problem, Rice’s Theorem
- complexity theory
 - P, NP, NP-completeness, Cook-Levin theorem

How can we make the theory course more accessible, more practical, and less intimidating?

Previous work has made substantial strides in this direction:

- interactive automata software tools such as JFLAP and DEM
 - Chesñevar et al (2003), Rodger et al (2006, 2009, etc.)
- “NP-completeness for all”
 - Crescenzi (2013), Enstrom (2010), Lobo (2006)
- Different theoretical model
 - Mandrioli (1982), Goldreich (2006, 2010)

Today's talk

How can we make the theory course more accessible, more practical, and less intimidating?


- Use *nondecision problems* as the primary paradigm
 - Contrasts with traditional decision problem paradigm

 Main emphasis of the paper

- Use *real computer programs* as the primary computational model
 - Contrasts with traditional use of Turing machines

 Also important (see the book)

A long-term vision for the theory course

- Make it more accessible and more practical
- Teach it to a wider range of undergraduates in more institutions
- Place it earlier in the curriculum with fewer prerequisites
- How?
 - Use nonddecision problems  Next: explain what this means
 - Use real computer programs

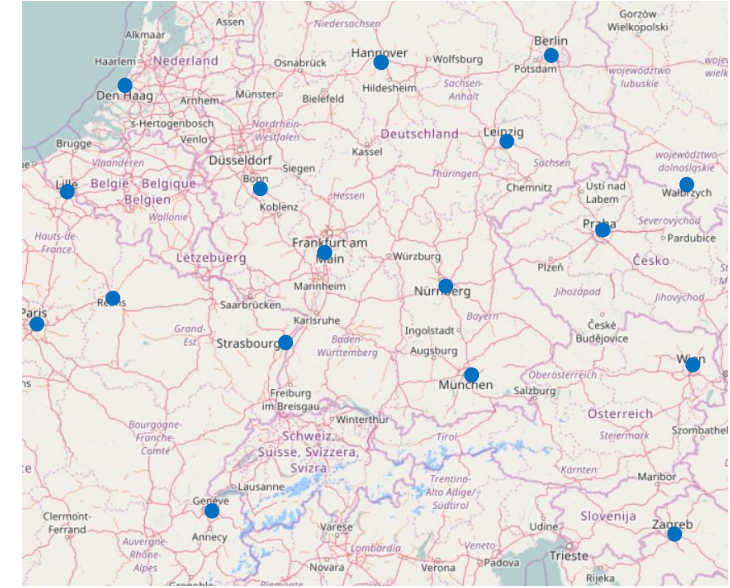
Which is more “useful”: program *A* or program *B*?

Input: Input to both programs is a roadmap and a list of cities:

Output:

Program *A* outputs { **“yes”** if there’s a driving route that visits each city and takes less than 100 hours
“no” otherwise

Program *B* outputs { **a description of a suitable route** if there’s a driving route that visits each city and takes less than 100 hours
“no” otherwise



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Which is more relevant for teaching: program *A* or program *B*?

Program *A* outputs { **“yes”**
“no”

Program *B* outputs { **a description of a
suitable route**
“no”

Which is more relevant for teaching: program *A* or program *B*?

Program *A* outputs { **“yes”** • *Decision problem.*
“no”

Program *B* outputs { **a description of a
suitable route** • *Nondecision problem.*
“no”

Which is more relevant for teaching: program *A* or program *B*?

Program *A* outputs { **“yes”**
“no”

- *Decision problem.*
- Existing theory-of-computation courses usually focus on decision problems.

Program *B* outputs { **a description of a
suitable route**
“no”

- *Nondecision problem.*

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- *Nondecision problem.*

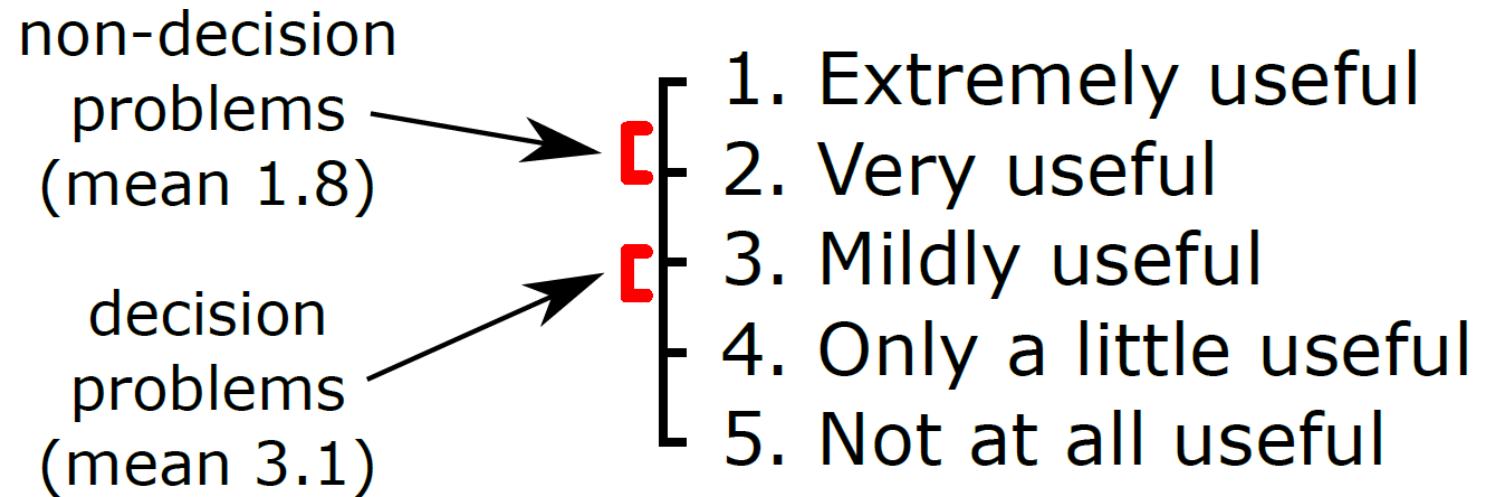
- This talk points to a way to teach the theory-of-computation course using *nondecision problems* and *real computer programs*.
- Students may achieve better learning because the content is perceived as relevant and practical.

We consider only a *novice audience*

- Novice audience \equiv undergraduate students who are seeing computability and complexity theory for the first time
- Experienced practitioners know that decision programs can often be converted to equivalent non-decision programs with only a logarithmic increase in running time.
- But for the *novice audience*, a program that outputs only a single bit may appear abstruse, irrelevant, and impractical

Programs that solve nondecision problems are perceived as much more “useful” by the novice audience

Survey of undergraduates compared decision and nondecision variants of TSP and knapsack problems



- The difference has overwhelming statistical significance and large effect size
- Perceived usefulness translates to better learning outcomes (Fink 2013)

OK, we need to use nonddecision problems.
But how do we do that?

- Answer: adjust certain definitions
- A brief example is shown next
- Please see the paper for details

Example of a technical detail: computational problems

- A **computational problem** (which may or may not be a decision problem) is a function F , mapping ASCII strings to sets of ASCII strings.
- If $F(x) = \{s_1, s_2, \dots\}$, we call $\{s_1, s_2, \dots\}$ the **solution set** for x , and each s_i is a **solution** for x .

This allows us to talk about “solving a problem” instead of “deciding a language”:

- Computer program P **solves** the computational problem F if $P(x) \in F(x)$ for all x .
- Contrast with: Turing machine M **decides** language L if M accepts all $s \in L$ and rejects all $s \notin L$

Examples of “solving a problem” instead of “deciding a language”

| | Traditional (decision) | Practical (nondecision) |
|-----------------|--|---|
| HamCycle | Does this graph have a Hamilton cycle? e.g. “a,b b,c c,a” \mapsto “yes” | Please give me a Hamilton cycle of this graph. e.g. “a,b b,c c,a” \mapsto “a,b,c” |
| Factor | Does this integer have a nontrivial factor? e.g. “51295697” \mapsto “yes” | Please give me a nontrivial factor of this integer. e.g. “51295697” \mapsto “8779” |

Other examples of technical details

- To incorporate nondcision problems, we need generalizations of traditional complexity classes, e.g.:
 - P becomes Poly
 - NP becomes NPoly
- The generalized definition of a *verifier* offers new pedagogical opportunities:
 - The traditional role of the *certificate* is separated into two clearer, independent notions: the *solution* and the *hint*

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Why? As Turing himself wrote (1936):

This proof, although perfectly sound, has the disadvantage that it may leave the reader with a feeling that “there must be something wrong”.

Using real computer programs permits interactive experimentation by students

Example 1: A classical diagonalization + proof by contradiction can be done explicitly in Python

```
from yesOnString import yesOnString
def weirdYesOnString(progString):
    if yesOnString(progString, progString)=='yes':
        return 'no'
    else:
        return 'yes'
```

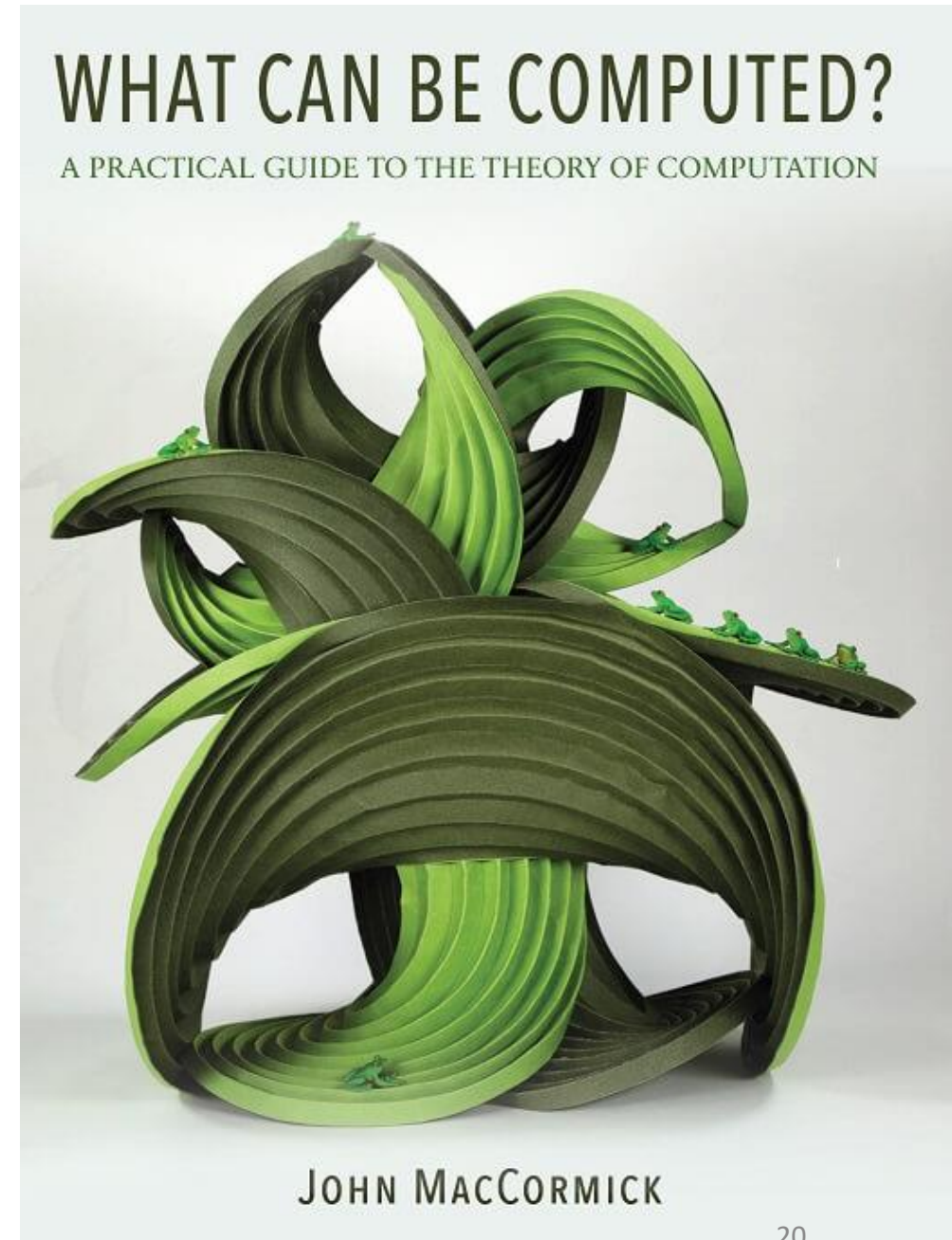
Using real computer programs permits interactive experimentation by students

Example 2: a “universal program” is much simpler than a universal Turing machine.

```
def universal(progString, inString):  
    # Execute the definition of the function in progString. This defines  
    # the function, but doesn't invoke it.  
    exec(progString)  
    # Now that the function is defined, we can extract a reference to it.  
    progFunction = utils.extractMainFunction(progString)  
    # Invoke the desired function with the desired input string.  
    return progFunction(inString)
```

That sounds interesting. But how can I actually teach a theory course using nonddecision problems and real computer programs?

- Answer: There is a new text book from Princeton University Press that takes this approach
 - *What Can Be Computed?: A Practical Guide to the Theory of Computation*
 - Available Spring 2018
 - Visit the Princeton University Press booth or email me for more details



Conclusion: a long-term vision for the theory course

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- Place it earlier in the curriculum with fewer prerequisites
- How?
 - Use nondcision problems
 - Use real computer programs
- Thanks for listening!

